

Efficient Computation of $\text{tr}\{\mathbf{TR}^{-1}\}$ for Toeplitz Matrices

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Abstract—An efficient algorithm for the computation of $\text{tr}\{\mathbf{TR}^{-1}\}$, where \mathbf{T} and \mathbf{R} are Toeplitz matrices and \mathbf{R} is also symmetric positive definite, is presented. The method exploits the fact that the trace of \mathbf{TR}^{-1} depends only on the sum of the diagonals of \mathbf{R}^{-1} , and not on the whole matrix \mathbf{R}^{-1} . To obtain this sum, a fast efficient technique, built upon the Trench algorithm for computing the inverse of a Toeplitz matrix, is developed. The complexity of the algorithm depends on the generation function of matrix \mathbf{R} and is $O(N \ln N)$ for generic functions and $O(p \ln p)$ for AR(p) functions.

Index Terms—Fast algorithm, fast Fourier transform, Toeplitz matrix, trace, Trench algorithm.

I. INTRODUCTION

THE NEED for computing $\text{tr}\{\mathbf{TR}^{-1}\}$, where \mathbf{T} and \mathbf{R} are Toeplitz matrices and \mathbf{R} is also symmetric and positive definite (SPD), appears in many signal processing problems. Relevant examples are estimation of Toeplitz constrained covariance matrices [1], [2], matrix approximation under the Frobenius norm [3], functional approximation of Gaussian densities using Kullback divergence [4], channel estimation [5], pulse time-of-arrival analysis [6], and computation of the Fisher information matrix of zero-mean Gaussian processes [7]. In the latter case, one has to compute terms of the form $\text{tr}\{\mathbf{T}\delta\mathbf{R}^{-1}\}$, where $\delta\mathbf{R}^{-1}$ denotes the derivative of \mathbf{R}^{-1} with respect to a given parameter. These terms can be determined numerically from the knowledge of $\text{tr}\{\mathbf{TR}^{-1}\}$.

Given the Toeplitz matrices $\mathbf{T} \equiv [t_{i-j}]$ and $\mathbf{R} \equiv [r_{i-j}]$, for $i, j = 1, \dots, N$, the obvious way of computing $\text{tr}\{\mathbf{TR}^{-1}\}$ is to invert \mathbf{R} and then determine the trace of \mathbf{TR}^{-1} ; by using the Trench algorithm (see, e.g., [8]) to compute \mathbf{R}^{-1} , the total complexity in computing $\text{tr}\{\mathbf{TR}^{-1}\}$, measured in floating point operations, is $(9/4)N^2 + (3/2)N$ (\mathbf{R}^{-1} takes $(7/4)N^2$ and the remaining operations take $N^2/2 + (3/2)N$).

Porat and Friedlander [9], based on the Levinson–Durbin algorithm for computing the orthogonal polynomials of a Toeplitz matrix, proposed an algorithm for the exact computation of the Fisher information matrix. This algorithm can be easily adapted to compute $\text{tr}\{\mathbf{TR}^{-1}\}$. The complexity of the method is still $O(N^2)$.

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When $N \rightarrow \infty$, and under adequate hypothesis on sequences $\{t_\tau, \tau \in \mathbb{Z}\}$ and $\{r_\tau, \tau \in \mathbb{Z}\}$ (\mathbb{Z} denotes the integer set), we have (see, e.g., [7, p. 140])

$$\lim_{N \rightarrow \infty} N^{-1} \text{tr}\{\mathbf{TR}^{-1}\} = \frac{1}{2\pi} \int_{-\pi}^{\pi} S_t(\omega) S_r^{-1}(\omega) d\omega \quad (1)$$

where $S_t(\omega)$ and $S_r(\omega)$ are the discrete Fourier transform of sequences $\{t_\tau, \tau \in \mathbb{Z}\}$ and $\{r_\tau, \tau \in \mathbb{Z}\}$, respectively. Asymptotical result (1) is the basis of Whittle’s formula [10], for the asymptotic normalized Fisher information matrix of a zero-mean normal process. This formula (1), despite leading to closed and simple expressions with light complexity, yields frequently a poor approximation for *small* sample sizes (see e.g., [11]).

To our knowledge, there is no general technique for the determination of $\text{tr}\{\mathbf{TR}^{-1}\}$ with complexity lower than $O(N^2)$. In this letter, we introduce a faster technique. We begin by noting that $\text{tr}\{\mathbf{TR}^{-1}\}$ depends only on the sum of the diagonals of \mathbf{R}^{-1} (which we refer to as the *diagonal sum* of \mathbf{R}^{-1}). With this fact in mind, it is then proved, based on the Trench algorithm for determining the inverse of a Toeplitz matrix, that the diagonal sum of the referred matrix can be computed with $O(N \ln N)$ complexity. In computing the diagonal sum, it is necessary to solve a Toeplitz system. By using the preconditioned conjugate gradient technique (see [12]–[15]), this step has $O(N \ln N)$ complexity. If matrix \mathbf{R} is generated by a rational function of order (p, q) , the methods [16], [17] solve the system with complexity $\max(p, q)O(N)$. Therefore, the total complexity in computing $\text{tr}\{\mathbf{TR}^{-1}\}$ does not exceed $O(N \ln N)$.

II. EFFICIENT COMPUTATION OF $\text{tr}\{\mathbf{TR}^{-1}\}$

Let $\mathbf{T}, \mathbf{R} \in \mathbb{R}^{N \times N}$ be Toeplitz matrices¹ of real elements, where \mathbf{R} is SPD and $\mathbf{R}^{-1} \equiv [c_{ij}]$. Noting that $t_{ij} = t_\tau$, where $\tau \equiv i - j$ for $i, j = 1, \dots, N$, it follows that

$$\begin{aligned} \text{tr}\{\mathbf{TR}^{-1}\} &= \sum_{i,j=1}^N t_{ij} c_{ji} = \sum_{\tau=-N+1}^{N-1} t_\tau \sum_{i \in S_\tau} c_{i-\tau, i} \\ &= \sum_{\tau=-N+1}^{N-1} t_\tau \bar{c}_{-\tau} \end{aligned} \quad (2)$$

¹Sometimes, we use the subscript N meaning that matrices \mathbf{T}_N and \mathbf{R}_N are of dimension $N \times N$.

where $\bar{c}_\tau \equiv \sum_{i \in S_\tau} c_{i+\tau, i}$ is the τ th diagonal sum of \mathbf{R}^{-1} ($\tau \geq 0$ denotes south-west diagonals), and

$$S_\tau = \begin{cases} 1 + \tau, \dots, N & \tau \geq 0 \\ 1, \dots, N - |\tau| & \tau < 0. \end{cases} \quad (3)$$

According to (2), $\text{tr}\{\mathbf{TR}^{-1}\}$ depends on t_τ and on \bar{c}_τ (sum of the elements of \mathbf{R}^{-1} along diagonal τ).

A. Sum of Diagonals of \mathbf{R}^{-1}

Toeplitz matrices belong to the larger class of *persymmetric matrices* [18]: matrix \mathbf{R}_N is persymmetric if it is symmetric about its northeast–southwest diagonal, i.e., if $r_{ij} = r_{N-j+1, N-i+1}$ for $i, j = 1, \dots, N$. In an equivalent form $\mathbf{R}_N = \mathbf{E}_N \mathbf{R}_N^T \mathbf{E}_N$, where $\mathbf{E}_N = [\delta_{N-i+1, j}]$ is the $N \times N$ exchange matrix. Note that $\mathbf{E}_N^{-1} = \mathbf{E}_N$. Thus, the inverse of a persymmetric matrix is, if it exists, also persymmetric.

Consider the partition

$$\mathbf{R}_N^{-1} = \begin{bmatrix} \mathbf{R}_{N-1} & \mathbf{E}_{N-1} \mathbf{r} \\ \mathbf{r}^T \mathbf{E}_{N-1} & r_0 \end{bmatrix}^{-1} = \begin{bmatrix} \mathbf{R}' & \boldsymbol{\nu} \\ \boldsymbol{\nu}^T & \gamma \end{bmatrix} \quad (4)$$

with $\mathbf{r} = [r_{-1}, \dots, r_{-N+1}]^T$. Vector $[\boldsymbol{\nu}^T, \gamma]^T$, the last column of \mathbf{R}_N^{-1} , and matrix \mathbf{R}' are given by (see, e.g., [18, p. 130])

$$\begin{cases} \boldsymbol{\nu} = \gamma \mathbf{E} \boldsymbol{\alpha} \\ \gamma = \frac{1}{r_0 + \mathbf{r}^T \boldsymbol{\alpha}} \\ \mathbf{R}' = \mathbf{R}_{N-1}^{-1} + \frac{\boldsymbol{\nu} \boldsymbol{\nu}^T}{\gamma} \end{cases} \quad (5)$$

where $\boldsymbol{\alpha}$ is the solution of the Yule–Walker equation $\mathbf{R}_{N-1} \boldsymbol{\alpha} = -\mathbf{r}$.

Matrix \mathbf{R}_{N-1}^{-1} exists and is persymmetric; thus $[\mathbf{R}_{N-1}^{-1}]_{i, j} = [\mathbf{R}_{N-1}^{-1}]_{N-j, N-i}$ and therefore

$$c_{ij} = [\mathbf{R}_{N-1}^{-1}]_{i, j} + \frac{1}{\gamma} \nu_i \nu_j \quad (6)$$

$$= [\mathbf{R}_{N-1}^{-1}]_{N-j, N-i} + \frac{1}{\gamma} \nu_i \nu_j \quad (7)$$

for $i, j = 1, \dots, N-1$. From (6) we see that $[\mathbf{R}_{N-1}^{-1}]_{N-j, N-i} = c_{N-j, N-i} - (1/\gamma) \nu_{N-j} \nu_{N-i}$. Using this equality in (7), we are led to

$$c_{ij} = c_{N-j, N-i} + \frac{1}{\gamma} (\nu_i \nu_j - \nu_{N-j} \nu_{N-i}). \quad (8)$$

Since $[c_{ij}] = \mathbf{R}^{-1}$ is persymmetric, element c_{ij} satisfies $c_{N-j, N-i} = c_{i+1, j+1}$ for $i, j = 1, \dots, N-1$; replacing $c_{N-j, N-i}$ by $c_{i+1, j+1}$ in (8), it follows that

$$c_{i+1, j+1} = c_{ij} + \frac{1}{\gamma} (\nu_{N-j} \nu_{N-i} - \nu_i \nu_j) \quad (9)$$

for $i, j = 1, \dots, N-1$. Defining $j \equiv i + \tau$, $\nu_N \equiv \gamma$, and noting that $c_{1, i} = \nu_{N+1-i}$ for $i, j = 1, \dots, N$, it follows that

$$c_{i+1, i+1+\tau} = c_{i, i+\tau} + \frac{1}{\gamma} (\nu_{N-i-\tau} \nu_{N-i} - \nu_i \nu_{i+\tau}) \quad (10)$$

TABLE I
SUMMARY OF THE PROPOSED SCHEME FOR THE DETERMINATION OF THE DIAGONAL SUM OF \mathbf{R}_N^{-1}

Computation of \bar{c}_τ	
Step	Complexity
1. Solve $\mathbf{R}_{N-1} \boldsymbol{\alpha} = -\mathbf{r}$	$O(N \ln N)$
2. Compute $[\boldsymbol{\nu}^T, \gamma]$ according (5)	N
3. Compute $\mathcal{DF}[u_i'' \nu_i']$ and $\mathcal{DF}[\nu_i']$	$2N \ln_2 2N$
4. Compute \bar{c}_τ according (14)	$N \ln_2 2N$

valid for $|\tau| \leq N-2$ and $i = 1, \dots, N-\tau-1$, when $\tau \geq 0$, and $i = 1-\tau, \dots, N-1$, when $\tau \leq 0$. Expressions (9) or (10) generate recursively, from vector $\boldsymbol{\nu}$, elements c_{ij} along each diagonal of \mathbf{R}^{-1} . The diagonal sum \bar{c}_τ is, after simple but lengthy manipulation of (10), given by

$$\gamma \bar{c}_\tau = \sum_{i=1}^{N-\tau} [i \nu_i \nu_{i+\tau} + \nu_i \nu_{i+\tau} (i + \tau) - N \nu_i \nu_{i+\tau}]. \quad (11)$$

The determination of \bar{c}_τ , for $\tau = 1, \dots, N-1$, according to (11), has $O(N^2)$ complexity. Notice, however, that each term in the sum (11) defines a convolution, which can be computed using the *fast Fourier transform* (FFT) with $O(N \ln N)$ complexity. For this purpose, define $\mathbf{0}_N^T$ as a zero row vector of dimension N and the sequences $\{u_i'\}$ and $\{\nu_i'\}$, with $i \in \mathbb{Z}$, as periodic extensions, of period $2N$, of sequences $\{1, \dots, N, \mathbf{0}_N^T\}$ and $\{\nu_1, \dots, \nu_N, \mathbf{0}_N^T\}$, respectively. Sum (11), using entities u_i' and ν_i' , assumes the form

$$\gamma \bar{c}_\tau = \sum_{i=1}^{2N} (u_i' \nu_i' \nu_{i+\tau}' + \nu_i' \nu_{i+\tau}' u_{i+\tau}' - N \nu_i' \nu_{i+\tau}') \quad (12)$$

$$= \underbrace{(u_{-i}' \nu_{-i}') \star \nu_{-i}' + \nu_{-i}' \star (u_i' \nu_i')}_{A_{N, \tau}} - N \underbrace{\nu_{-i}' \star \nu_i'}_{B_{N, \tau}} \quad (13)$$

where symbol \star means circular convolution of length $2N$. Denoting $\mathcal{DF}[x]$ as the $2N$ -vector containing the discrete time Fourier series of x and \mathcal{DF}^{-1} its inverse (i.e., $x_n = \mathcal{DF}^{-1}[\mathcal{DF}[x]]_{(n)}$), and using the discrete time Fourier series properties (convolution, time symmetry, and conjugation), expression (13) is given by

$$\gamma \bar{c}_\tau = \mathcal{DF}^{-1} \{ \text{Re}(\mathcal{DF}[(u_i'' \nu_i')] \odot \mathcal{DF}^*[\nu_i']) \}_{(-\tau)} \quad (14)$$

where $u_i'' \equiv 2u_i' - N$ and symbol \odot denotes element-wise multiplication.

The number of floating point operations needed to implement (14) is, approximately, $3N \ln 2N$ (corresponding to three FFTs of size $2N$). Table I summarizes the proposed algorithm.

In the next section, we show that, if \mathbf{R} is the covariance matrix of an AR(p) process, then the terms γ , $A_{N, \tau}$ and $B_{N, \tau}$ of expression (13) are constant for $N > p$.

III. AUTOREGRESSIVE PROCESSES

Suppose that \mathbf{R} is a covariance matrix of an AR(p) process with coefficients $\{a_k, 0 \leq k \leq p\}$ where $a_0 \equiv 1$ and $a_p \neq 0$,

By hypothesis, filter $a(z) = \sum_{k=0}^p a_k z^{-k}$ is stable [i.e., $a(z)$ does not have zeros for $|z| \geq 1$]. In these conditions, we have for $p \geq N + 1$ (see, e.g., [7, ch. 6])

$$\gamma_{N+1} = \gamma_N \quad (15)$$

$$\boldsymbol{\alpha}_N = (a_1, a_2, \dots, a_p, 0, \dots, 0). \quad (16)$$

From the expression for $\boldsymbol{\nu}$ given in (5), we conclude that, for $N > p$

$$(\nu_1, \dots, \nu_{N-1}, \nu_N) = \gamma(0, \dots, 0, a_p, \dots, a_1, a_0). \quad (17)$$

Replacing (17) into (13), assuming that $N > p$, and after some manipulation, we get

$$\gamma \bar{c}_\tau = - \underbrace{\left(\sum_{i=0}^{p-\tau} i a_i a_{i+\tau} + \sum_{i=0}^{p-\tau} a_i a_{i+\tau} (i + \tau) \right)}_{A_{N,\tau}} + N \underbrace{\left(\sum_{i=0}^{p-\tau} a_i a_{i+\tau} \right)}_{B_{N,\tau}}. \quad (18)$$

We conclude then that terms γ , $A_{N,\tau}$ and $B_{N,\tau}$ do not depend on N for $N > p$. Therefore, the computation of $\text{tr}\{\mathbf{TR}^{-1}\}$, where \mathbf{R} is the covariance matrix of an AR(p) process, according to the proposed method, has $O(p \ln p)$ complexity for $N > p$.

For generic processes, formula (18) does not apply. However, terms γ_N , $A_{N,\tau}$, and $B_{N,\tau}$ converge to a constant as values of $N \rightarrow \infty$. This is a consequence of the Levinson–Durbin recursions

$$\boldsymbol{\alpha}_N = [\boldsymbol{\alpha}_{N-1}, 0] - c_N [0, \mathbf{E}_N \boldsymbol{\alpha}_N]$$

where c_N is the N th-order reflection coefficient, which satisfy (see, e.g., [9])

$$\lim_{N \rightarrow \infty} \sum_{i=N+1}^{\infty} c_i^2 = 0.$$

Hence, for N_0 sufficiently large we have, for $N \geq N_0$

$$c_\tau \simeq \gamma_{N_0}^{-1} (A_{N_0,\tau} + N B_{N_0,\tau}).$$

The problem in applying the above extrapolation is, of course, how to determine N_0 , in order to obtain a good approximation for c_τ . A rule of thumb is to determine γ_N , $A_{N,\tau}$, and $B_{N,\tau}$ for an increasing sequence N_1, N_2, \dots , for example $N_i = 2^i$, for $i = 1, 2, \dots$, and estimate an integer N_i above which the former terms are practically constant.

IV. CONCLUDING REMARKS

We presented an efficient algorithm for computing $\text{tr}\{\mathbf{TR}^{-1}\}$, where \mathbf{T} and \mathbf{R} are Toeplitz matrices and \mathbf{R} is also symmetric and positive definite. The complexity of the algorithm depends on the generation function of matrix \mathbf{R} and is $O(N \ln N)$ for generic functions and $O(p \ln p)$ for AR(p) functions. We have also presented an approximated extrapolation formula with $O(N_0 \ln N_0)$ complexity, thus independent of N .

The central idea exploited is that $\text{tr}\{\mathbf{TR}^{-1}\}$ depends only on the sum of diagonals of \mathbf{R}^{-1} and not on each single element of this matrix. The computation of the referred sum is carried out efficiently by using the fast Fourier transform.

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